

Refinement Types for Stainless

Katarzyna Marek, Samuel Chassot, Matt Bovel

Outline

1. Dotty_(fork) has refinement types!

But a limited solver.

2. Stainless has a powerful solver!

But no refinements types.



3. Let's integrate Dotty_(fork) with
Stainless!

Refinement Types in Dotty_(fork)

Refinement Types

Refinement types are types qualified with logical predicates. For example,

`{x: Int with x > 0}`

denotes the type of all integers x such that $x > 0$.

Implemented in many languages (Liquid Haskell, Boolean refinement types in F*, Subset types in Dafny, etc.).

Prior art in Scala: SMT-based checking of predicate-qualified types for Scala (Schmid and Kunčák, Scala Symposium 2016), Refined, Iron.

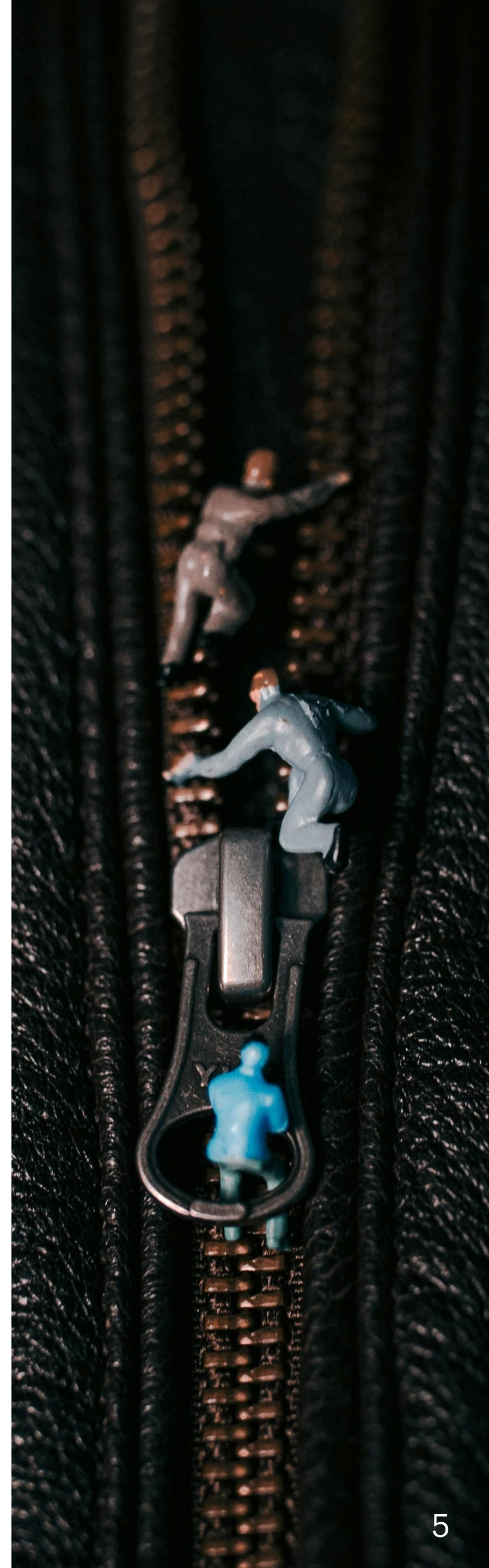
Example: Safe Lists Zipping

Consider the zip function:

```
def zip[A, B](xs: List[A], ys: List[B]): List[(A, B)]
```

Scala doesn't require both lists to have the same size:

```
val pairs = List(1, 2).zip(List("a"))  
// List((1, "a"))
```



Example: Safe Lists Zipping

Requiring same sizes,

```
def zip[A, B](
  xs: List[A],
  ys: List[B] with ys.size == xs.size,
): {res: List[A, B] with res.size = xs.size}
```

Usage:

```
val l1: List[Int] with l1.size == 3 = ...
val l2: List[String] with l2.size == 3 = ...
zip(l1, l2): List[(Int, String)] with l3.size == 3
```

Need to show:

```
{l2: List[String] with l2.size == 3}
<: {l2: List[String] with l2.size == l1.size}
```



Limited Solver

Refinement types in Dotty ships with a minimal solver. It can:

- Reason about equality (like on the previous slide)
- Constant-fold and normalize: $1 + 2 \leftrightarrow 3$ and $y + z \leftrightarrow z + y$
- Order: $x > 2 \ \&\& \ y > x \rightarrow y > 2$, but not $x < 10 \rightarrow x + 2 < 8$
- Unfolding:

```
def fact(n: Int):  
  {r: Int with r == (if n == 0 then 1 else n * fact(n - 1))}
```

$\text{fact}(1) = 1$, but not $\text{fact}(2) = 2$

Stainless Verifier

Stainless: Automated Proof

Verification Framework for Scala

```
def zip[A, B](xs: List[A], ys: List[B]): List[(A, B)] = {  
  require(xs.size == ys.size)  
  (xs, ys) match  
    case (Cons(x, xs0), Cons(y, ys0)) =>  
      Cons((x, y), zip(xs0, ys0))  
  
    case _ => Nil()  
}.ensuring (res => res.size == xs.size && res.map[A](p => p._1) == xs)
```

stainless summary

```
zip.scala:5:5: [Warning] Found counter-example: valid U:smt-cvc5 0.0  
zip.scala:9:11: zip postcondition valid U:smt-cvc5 0.1  
zip.scala:9:24: [Warning] xs: List[A] -> Cons[A](A#0, Nil[A]()) valid U:smt-cvc5 0.0  
zip.scala:9:24: zip precondition (call zip((scrut._1.t): @DropVCs, (scrut._...)) valid U:smt-cvc5 0.1  
zip.scala:11:15: [Warning] ys: List[B] -> Nil[B]() valid U:smt-cvc5 0.0  
-----  
total: 5 valid: 5 (0 from cache, 0 trivial) invalid: 0 unknown: 0 time: 0.29
```

Supported Scala Subset

Algebraic Data Types

Pure functional

Imperative variables

Local classes

While loops

Pattern matching

Type classes

Traits

Inner functions

Polymorphism

Type parameters

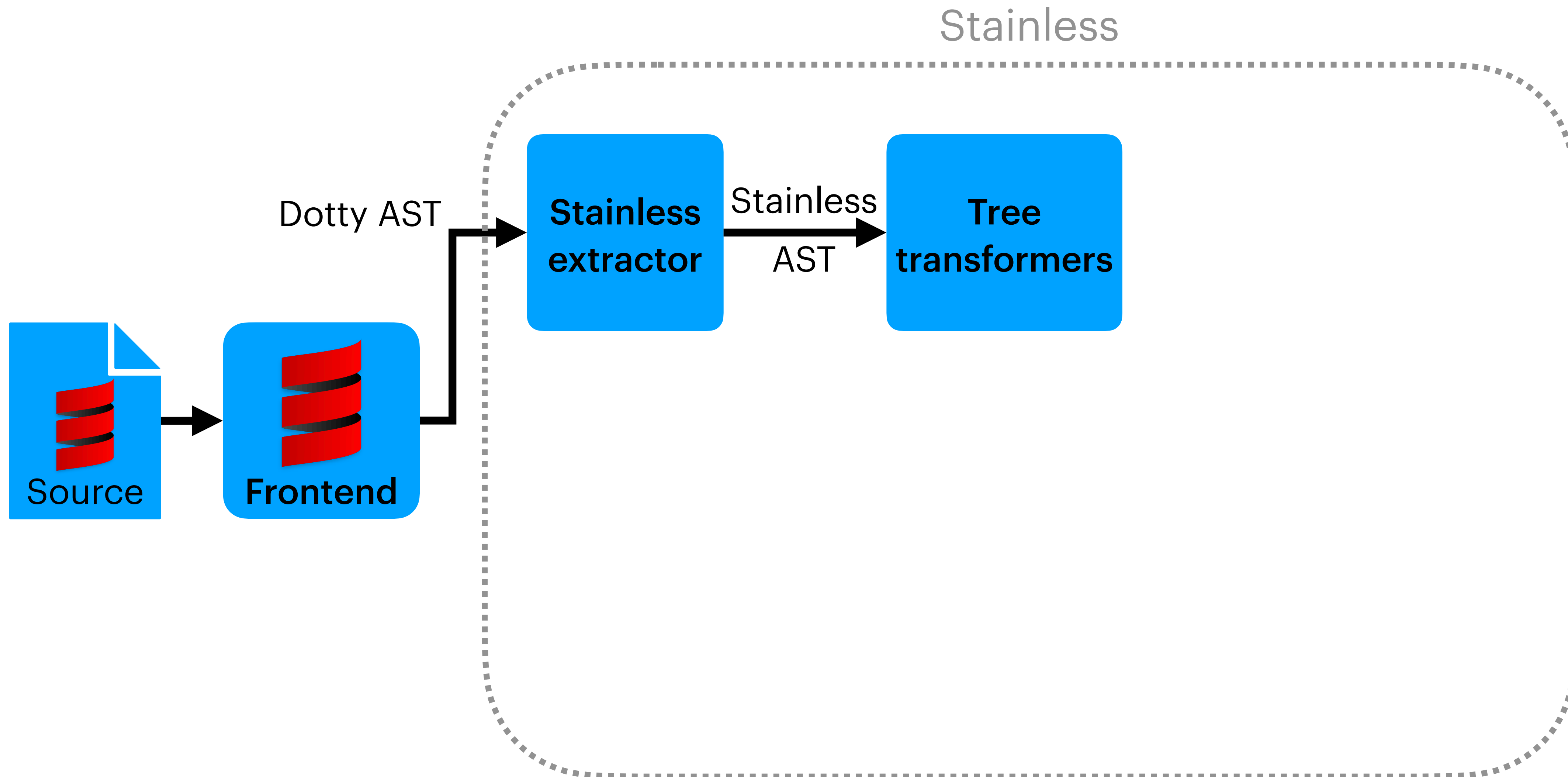
(Unreachable) exceptions

Type members

Type alias

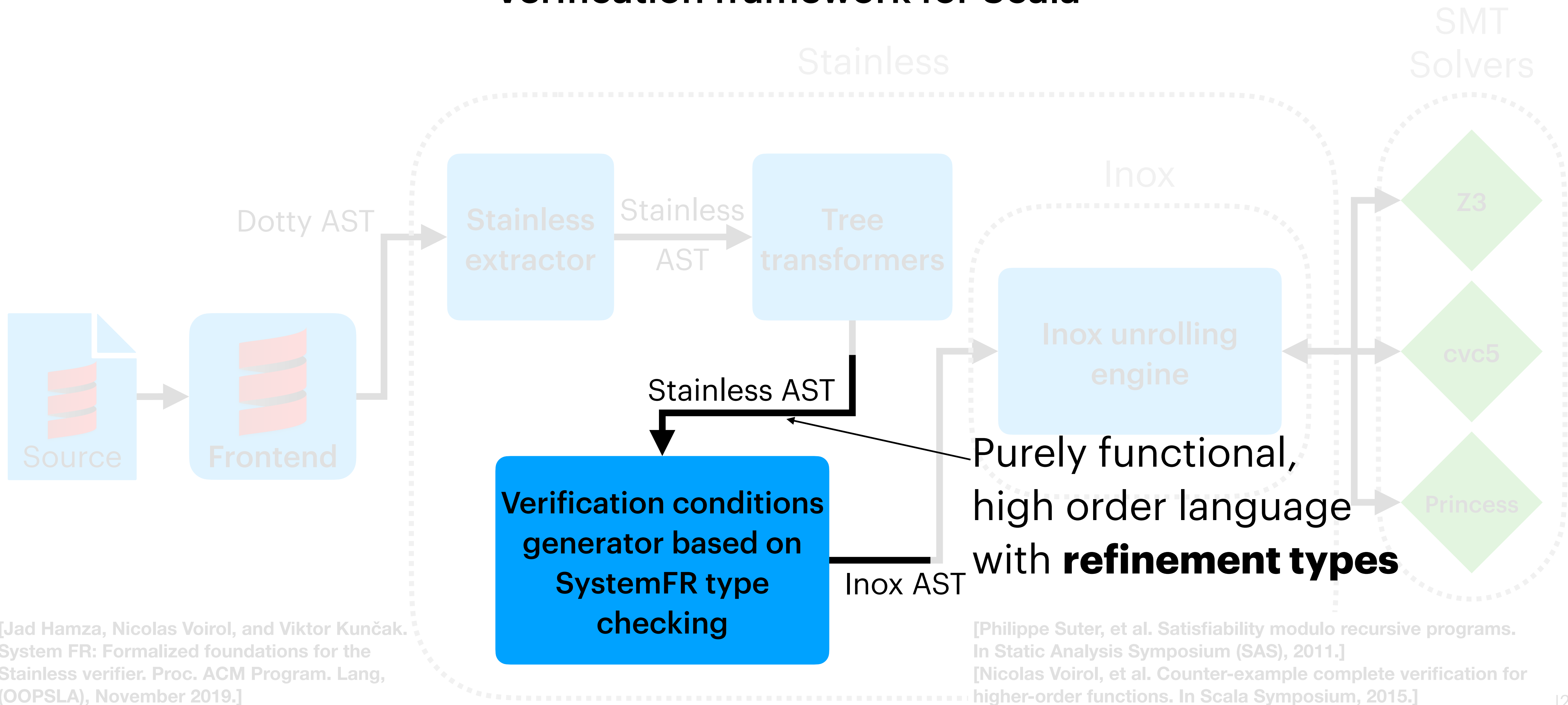
Stainless

Verification framework for Scala



Stainless

Verification framework for Scala



Internal trees

Before VC generation phase

```
def zip(xs: Cons[Int], ys: Cons[Boolean]): Cons[(Int, Boolean)] = {  
  require(xs.size == ys.size)  
  decreases(xs.size)  
  (xs.t, ys.t) match  
    case (xs0 @ Cons(_, _), ys0 @ Cons(_, _)) => Cons((xs.h, ys.h), zip(xs0, ys0))  
    case _ => Cons((xs.h, ys.h), Nil())  
}.ensuring(res => res.size == xs.size)
```

Internal trees

Before VC generation phase

```
def zip(xs: { v: List[Int] | v.isInstanceOf[Cons] },
        ys: { v: List[Boolean] | v.isInstanceOf[Cons] }
        ): { v: List[(Int, Boolean)] | v.isInstanceOf[Cons] } = {
  require(size[Int](xs) <= size[Boolean](ys))
  decreases(size[Int](xs))
  val targetBound: (List[Int], List[Boolean]) = (xs.t, ys.t)
  targetBound match {
    case (xs0 @ Cons(_, _), ys0 @ Cons(_, _)) =>
      Cons[(Int, Boolean)]((xs.h, ys.h), zip(xs0, ys0))
    case _ =>
      Cons[(Int, Boolean)]((xs.h, ys.h), Nil[(Int, Boolean)]())
  }
}.ensuring {
  (res: { v: List[(Int, Boolean)] | v.isInstanceOf[Cons] }) =>
    size[(Int, Boolean)](res) == size[Int](xs)
}
```

Stainless internal trees already rely on refinement types

Refinement Types + Stainless Verifier

Zip

```
def zip[A, B](a: List[A], b: List[B] with b.size == a.size ):  
  { l: List[(A, B)] with l.size == a.size  
    && l.map[A](_. _1) == a  
    && l.map[B](_. _2) == b } =  
  
(a, b) match  
  case (Nil(), Nil()) => Nil[(A, B)]()  
  case (Cons(h1, t1), Cons(h2, t2)) =>  
    Cons((h1, h2), zip(t1, t2))
```

stainless summary

```
./../Zip.scala:4:5: zip non-negative measure valid from cache 0.0  
./../Zip.scala:8:3: zip body assertion: match exhaustiveness valid from cache 0.0  
./../Zip.scala:8:3: zip type-checking valid U:smt-z3 0.1  
./../Zip.scala:9:28: zip cast correctness valid U:smt-z3 0.1  
./../Zip.scala:9:28: zip type-checking valid U:smt-z3 0.1  
./../Zip.scala:11:7: zip cast correctness valid U:smt-z3 0.2  
./../Zip.scala:11:7: zip type-checking valid U:smt-z3 0.0  
./../Zip.scala:11:22: zip measure decreases valid from cache 0.0  
./../Zip.scala:11:30: zip cast correctness valid from cache 0.0  
./../Zip.scala:11:30: zip precondition. (argument types: call zip[A, B]((scrut._1.t): @DropVCs , { ...}) valid from cache 0.0  
-----  
total: 10 valid: 10 (5 from cache, 0 trivial) invalid: 0 unknown: 0 time: 0.59
```

Scala Refinement Types + Stainless Verifier

Powerful solver for refinements
(for a subset of Scala),

- e.g., verifies the zip example

Additional features:

- opt-out termination checking
- advanced exhaustivity checks

Stainless Verifier + Refinement Types

Specification in the signature

New syntax?

Interface

```
trait Sorting:  
  def sort[T : Ordering](xs: List[T]): List[T] = {  
    (??? : List[T])th l.isSorted }  
  }.ensuring(res => res.isSorted)
```

```
class MergeSort extends Sorting { ... }
```

```
class InsertionSort extends Sorting { ... }
```

Higher-order functions

```
case class Pos(i: Int):  
  require(i > 0)
```

```
def toRoots(xs: List[Pos]): List[Int] =  
  xs.map(e => squareRoot(e.i))  
  xs.map(squareRoot)
```

```
def squareRoot(i: Int): Int =  
  require(i > 0)  
  require(i > 0)  
  ???  
  ???
```

stainless summary stainless summary

./HigherOrderExampleHigherOrderExampleRoots:1:10: toRoots depends on call to squareRoot(i) validation has failed -z3 0.0

total: 1 valid total: 1(1 from cache, 0 from local, 0 from trivial) invalid 1 time known 0.04 time: 0.34

Mutation

Intuitive semantics

```
def f(...): {res with T: p(res)}
```

Refinements on result types of **functions** are **postconditions**

```
val x: T with p(x)
```

Refinements on **values** are **assertions**

```
var x: T with p(x); sort[T : Ordering](xs: List[T])
```

Refinements on **variables** are

```
val ls: List[Int] with ls.isSorted =  
  List(1, 2, 4)
```

invariants

```
var ls: List[Int] with ls.isSorted = List(1, 2, 4)  
  List(1, 2, 4)  
  ls = List(5, 4) //error  
ls = List(5, 4) //error
```

```
val ls0: List[Int]  
  with ls0.isSorted =  
  List(1, 2, 4)
```

```
val ls1: List[Int]  
  with ls1.isSorted =  
  List(5, 4) //error
```

References to mutable state

```
case class Box(var v: Int):  
  def setV(y: Int):  
    {r: Unit with v == y} =  
      v = y
```

```
val b = Box(3)  
val p1: Unit with b.v0 == 1  
  = b.setV(1)  
val p2: Unit with b.v1 == 5  
  = b.setV(5)
```

```
var p1: Unit with b.v0 == 1  
  = b.setV(1)  
val p2: Unit with b.v1 == 5  
  = b.setV(5)
```

```
(p1: Unit with b.v0 == 1) = p1 // ok  
(p1: Unit with b.v1 == 1) = p1 // error
```

Refinements on result types of **functions** are **postconditions**

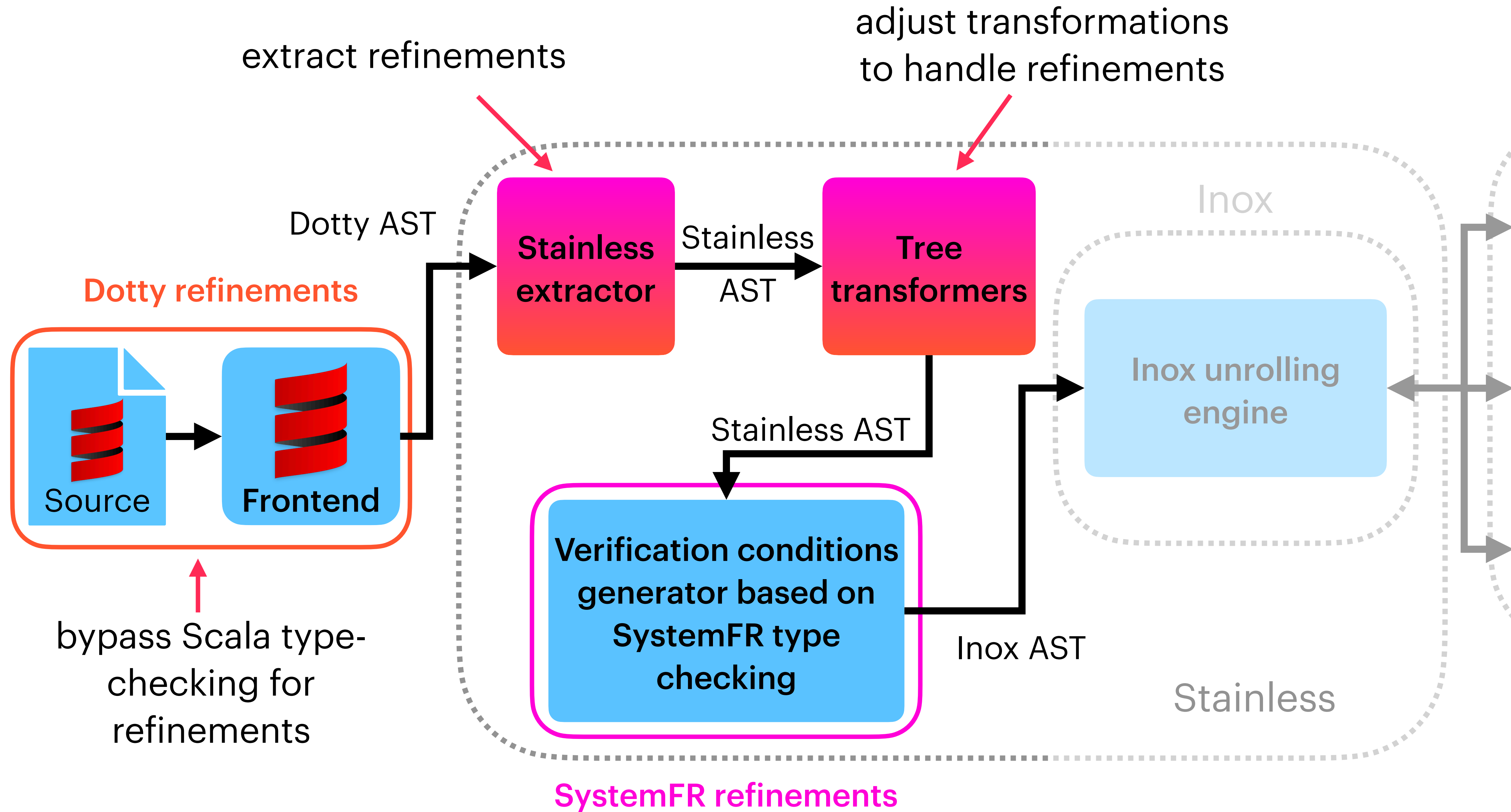
Refinements on **values** are **assertions**

```
val b0: Box̄ = Box(3)t  
val res1: () : Box, y0: Int):  
  (b0: Box, r: Unit with b0.v == 1)  
  = setV(b0, 1) b1.v == y0)  
val p1 = res1._2  
val b1 = res1._1  
val res2: () : variables is disallowed  
  (b1: Box, r: Unit with b1.v == 5)  
  = setV(b1, 5)
```

Intuitive semantics

	Pure refinements	References to mutable states	Side-effects in refinements
<code>def f(...): {res with T: p(res)}</code>	assertion/invariant	assertion	✗
<code>val x: T with p(x)</code>	assertion/invariant	assertion	✗
<code>var x: T with p(x)</code>	invariant	✗	✗

Implementation



Stainless
+
Refinement Types
=
**Powerful Verifier with
Expressive Specification**

- github.com/epfl-lara/dotty
- github.com/epfl-lara/stainless
- github.com/epfl-lara/inox

Thank you. Questions?